

Schemes in Lean

Ramon Fernández Mir

June 24, 2019

Formalising mathematics

Idea: Use a computer to check every step in a proof.

Why?

- There are many errors in published work.
- Some proofs are long and tedious to check 'by hand'.
- Build a database of all mathematical knowledge.

1967

1984 1986

1992

1998

2008

2013 2015

- Automath, de Bruijn.
- Coq, Coquand and Huet.
- Isabelle, Paulson.
- HOL Light, Harrison.
- First proof the Kepler conjecture, Hales.
- Proof of the four colour theorem in Coq, Gonthier.
- Proof of the odd order theorem in Coq, Gonthier.
- Proof of the Kepler conjecture in Isabelle and HOL Light, Hales.



- Automath, de Bruijn.
- **Coq, Coquand and Huet.**
- Isabelle, Paulson.
- HOL Light, Harrison.
- First proof the Kepler conjecture, Hales.
- Proof of the four colour theorem in Coq, Gonthier.
- Proof of the odd order theorem in Coq, Gonthier.
- Proof of the Kepler conjecture in Isabelle and HOL Light, Hales.



- Automath, de Bruijn.
- Coq, Coquand and Huet.
- Isabelle, Paulson.
- HOL Light, Harrison.
- First proof the Kepler conjecture, Hales.
- Proof of the four colour theorem in Coq, Gonthier.
- Proof of the odd order theorem in Coq, Gonthier.
- Proof of the Kepler conjecture in Isabelle and HOL Light, Hales.



- Automath, de Bruijn.
- Coq, Coquand and Huet.
- Isabelle, Paulson.
- **HOL Light, Harrison.**
- First proof the Kepler conjecture, Hales.
- Proof of the four colour theorem in Coq, Gonthier.
- Proof of the odd order theorem in Coq, Gonthier.
- Proof of the Kepler conjecture in Isabelle and HOL Light, Hales.



- Automath, de Bruijn.
- Coq, Coquand and Huet.
- Isabelle, Paulson.
- HOL Light, Harrison.
- **First proof the Kepler conjecture, Hales.**
- Proof of the four colour theorem in Coq, Gonthier.
- Proof of the odd order theorem in Coq, Gonthier.
- **Proof of the Kepler conjecture in Isabelle and HOL Light, Hales.**



- Automath, de Bruijn.
- Coq, Coquand and Huet.
- Isabelle, Paulson.
- HOL Light, Harrison.
- First proof the Kepler conjecture, Hales.
- Proof of the four colour theorem in Coq, Gonthier.
- Proof of the odd order theorem in Coq, Gonthier.
- Proof of the Kepler conjecture in Isabelle and HOL Light, Hales.

Lean

Lean was launched in 2013 and is developed mainly by Leonardo de Moura at Microsoft Research. See <http://leanprover.github.io>.

Features:

- Based on the Calculus of Inductive Constructions.
 - ▶ Types: dependent types, inductive types, `Prop` and `Type` u.
 - ▶ Terms: types, constants, lambda abstractions and applications.
 - ▶ To give a proof of a statement $P : \text{Prop}$ we build a term $p : P$.
- Supports powerful automation.
- Still at an early stage but promising standard library: `mathlib`.

(Xena project: Lean sessions every Thursday at the MLC)

Demo

Schemes

In 1960, the first volume of *Éléments de géométrie algébrique*, written by Alexander Grothendieck, is published.

Goal: Establish new foundations of algebraic geometry.

Approach: Define a *scheme* and build upon it.

Example

Suppose we want to study the polynomials $x = 0$ and $x^2 = 0$ over \mathbb{C} .

- Varieties: $Z(x) = Z(x^2) = \{0\}$.
- Schemes: $\text{Spec}(\mathbb{C}[x]/(x)) = \{(0)\}$ and $\text{Spec}(\mathbb{C}[x]/(x^2)) = \{(x)\}$.

Idea of an *affine scheme*:

- The set of functions on a space forms a ring.
- Start with a ring R .
- Define a topological space $\text{Spec}(R)$.
- Attach to it a *sheaf* of functions $\mathcal{O}_{\text{Spec}(R)}$ such that $\mathcal{O}_{\text{Spec}(R)}(\text{Spec}(R)) = R$ but do it in a way that it also makes sense to consider subsets of $\text{Spec}(R)$.

If we do it right, we have:

$$\{\text{Rings}\} \Leftrightarrow \{\text{Affine schemes}\}$$
$$\text{Algebra} \Leftrightarrow \text{Geometry}$$

This project

In March 2018, Kevin, Chris and Kenny defined a scheme in Lean.

Problems:

- Chaotic repository.
- Long and unnecessarily technical proofs.
- The definition was only mathematically equivalent.

Objective: start from scratch and do it well.

Canonical isomorphisms

In *Étale Cohomology*, a book about schemes by J. S. Milne we find:

An injection is denoted by \hookrightarrow , a surjection by \twoheadrightarrow , an isomorphism by \cong , a quasi-isomorphism (or homotopy) by \simeq , and a canonical isomorphism by $\stackrel{df}{=}$. The symbol $X \stackrel{df}{=} Y$ means X is defined to be Y , or that X equals Y by definition.

It brings up some questions:

- What is a canonical isomorphism?
- How does it look like in Lean?
- Is it the same as equality?

Canonical isomorphisms

In *Étale Cohomology*, a book about schemes by J. S. Milne we find:

An injection is denoted by \hookrightarrow , a surjection by \twoheadrightarrow , an isomorphism by \cong , a quasi-isomorphism (or homotopy) by \simeq , and a canonical isomorphism by $\stackrel{df}{=}$. The symbol $X \stackrel{df}{=} Y$ means X is defined to be Y , or that X equals Y by definition.

It brings up some questions:

- What is a canonical isomorphism? An 'obvious' isomorphism.
- How does it look like in Lean? Inverse maps in both directions.
- Is it the same as equality? No!

Definition

Suppose R is a ring and $f \in R$ then the *localisation of R away from f* is $R_f = \{(a, f^n) \mid a \in R, n \in \mathbb{N}\} / \sim$ where $(a, f^n) \sim (b, f^m)$ if and only if $f^k(a f^m - b f^n) = 0$ for some k . We write $\frac{a}{f^n}$ to mean $[(a, f^n)]$.

We wanted to use a lemma about the sequence:

$$0 \longrightarrow R \xrightarrow{\alpha} \bigoplus_i^n R_{f_i} \xrightarrow{\beta} \bigoplus_{i,j}^n R_{f_i f_j}$$

But we didn't have R_{f_i} and $R_{f_i f_j}$. Instead, we had $A_i \cong R_{f_i}$ and $B_{ij} \cong R_{f_i f_j}$, so the situation was:

$$\begin{array}{ccccccc} 0 & \longrightarrow & R & \xrightarrow{\alpha} & \bigoplus_i^n R_{f_i} & \xrightarrow{\beta} & \bigoplus_{i,j}^n R_{f_i f_j} \\ & & \searrow \alpha' & & \psi_i \updownarrow \psi_i^{-1} & & \psi_{ij} \updownarrow \psi_{ij}^{-1} \\ & & & & \bigoplus_i^n A_i & \xrightarrow{\beta'} & \bigoplus_{i,j}^n B_{ij} \end{array}$$

Solution:

- Neil Strickland came up with three axioms that characterise the class of rings isomorphic to R_{f_i} (respectively $R_{f_i f_j}$).
- Prove the lemma about the sequence for any type satisfying the three axioms.
- Show that A_i (respectively B_{ij}) satisfies the axioms.
- Avoid arguing about α' , β' , ψ_i , ψ_i^{-1} , ψ_{ij} and ψ_{ij}^{-1} .

Locally ringed spaces

The definition of a locally ringed space has three components:

- 1 A topological space X .
- 2 A sheaf of rings on X , denoted \mathcal{O}_X .
- 3 The property that, for all $x \in X$, the *stalk* $\mathcal{O}_{X,x}$ is a *local ring*.

The previous definition included 1 and 2. 3 followed by construction.

Theorem

The topological space $\text{Spec}(R)$ has the structure of a locally ringed space.

Over 7000 lines of code building up to the following.

Definition

A *scheme* is a locally ringed space (X, \mathcal{O}_X) with an open cover $\{U_i\}$ such that for all i , $(U_i, \mathcal{O}_X|_{U_i})$ is isomorphic to $(\text{Spec}(R), \mathcal{O}_{\text{Spec}(R)})$ for some ring R .

```
structure scheme (X : Type u) [topological_space X] :=
  (carrier      : locally_ringed_space X)
  (Haffinecov  : ∃ (OC : covering.univ X), ∀ i,
    ∃ (R : Type v) [comm_ring R]
      (fpU : open_immersion_pullback (Spec R) carrier.O.F),
      fpU.range = OC.Uis i
      ∧ fpU.carrier ≅ structure_sheaf.presheaf R)
```

Conclusion and future work

We have given the first complete formal definition of a scheme and it is available in <https://github.com/ramonfmir/lean-scheme>.

Next steps:

- Kenny proved the correspondence between the category of rings and the category of affine schemes.
- Have an example of a projective scheme.
- Add schemes to mathlib!

Questions